Performance Evaluation of OpenCL Standard Support (and Beyond)

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Background

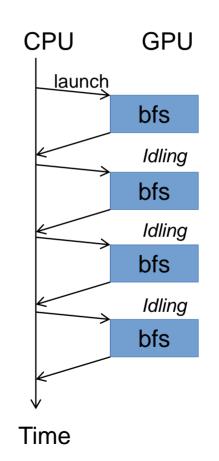
- IrGL: DSL for (Irregular) Graph Analytic Kernels
 - Wrote a compiler from IrGL to CUDA
 - Fastest graph kernels
 - Pai and Pingali, "A compiler for throughput optimization of graph algorithms on GPUs", OOPSLA 2016
- •Teamed up with Tyler and Ally to target OpenCL

IrGL Key Insight

- Graph algorithms suffer 3 bottlenecks
- Need 3 key optimizations for high performance
 - Iteration Outlining
 - Nested Parallelism (not OpenCL NP)
 - Cooperative Conversion

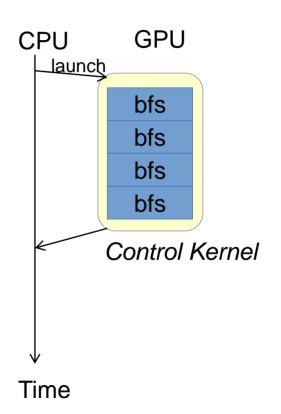
Problem: CPU—GPU Launch

- Most graph algorithms are iterative
 - Repeat until fixpoint
- If time per iteration is small (average ~20us for BFS), launch throughput can't keep up



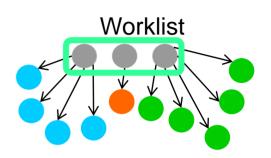
Optimization #1: Iteration Outlining

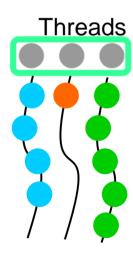
- •Compiler generates a control kernel that "launches" child kernels
 - Actually inlines them
- Need a global barrier between "kernel invocations" (now function calls)



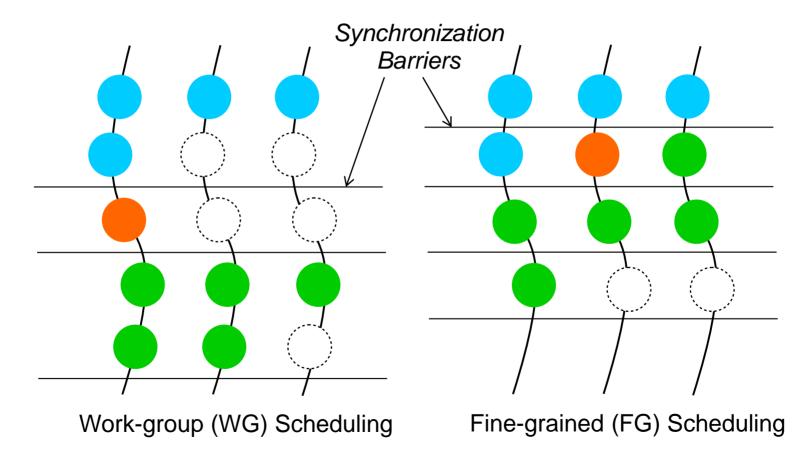
Problem: Load Imbalance

- Graph algorithms usually two *parallel* nested loops
 - Outer loop over nodes
 - Inner loop over edges of a node
- Graph edge distribution

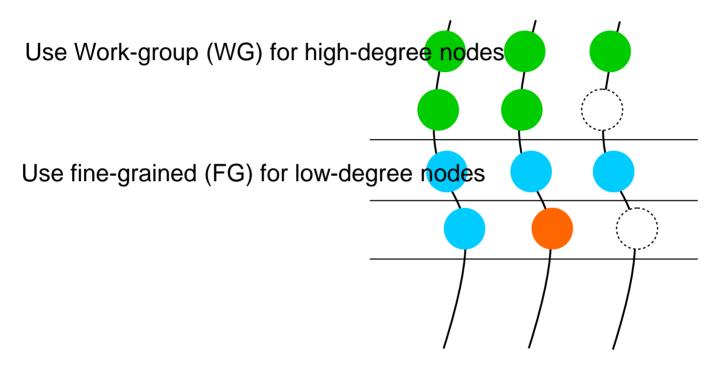




Optimization #2: Nested Parallelism



NP: Multiple Schedulers

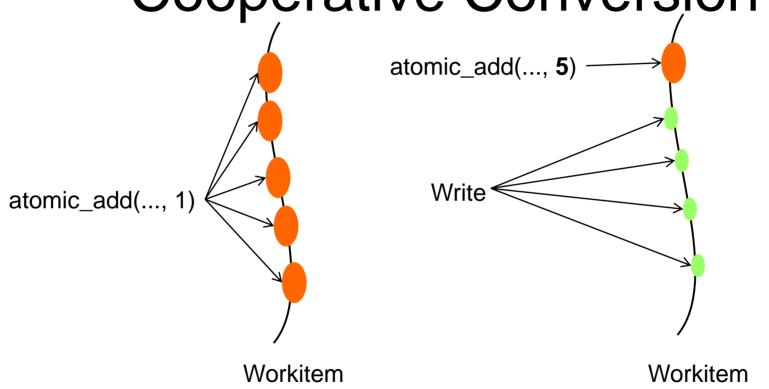


Work-group (WG) + Finegrained (FG) Scheduling

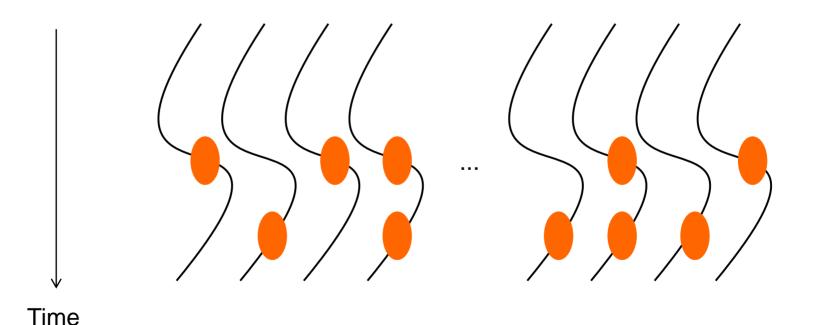
Implementing IrGL NP in OCL

- Support any combination of three schedulers
 - workgroup (wg)
 - subgroup (sg)
 - finegrained (fg)
- •Workgroup and Finegrained schedulers require:
 - local memory and workgroup barriers
- •Subgroup scheduler requires:

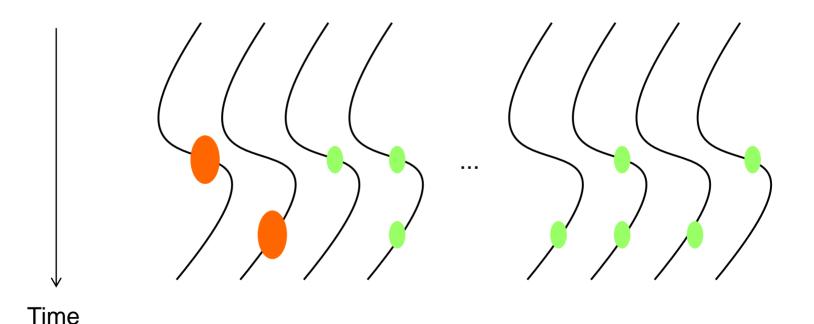
Optimization #3: Cooperative Conversion



Dynamic aggregation problem



Aggregate across workitems



Implementing Coop-Conv in OCL

- Supports aggregation:
 - within a workitem
 - across a subgroup
 - across all workitems of a workgroup
- •Workitems use prefix scans for aggregation, requiring barriers
 - Unlike CUDA, even subgroup aggregation

GPUs We Used

| Vendor | GPU | OpenCL version |
|--------|--------------|----------------|
| ARM | Mali 4 | 1.2 |
| NVIDIA | GTX 1080 | 1.2 |
| | Quadro M4000 | 1.2 |
| AMD | R9 | 2.0 |
| Intel | HD 5500 | 2.0 |
| | Iris 6100 | 2.0 |

Portable C11-style Atomics (OCL2)

- •Required for: coop-conv, wg, and oitergb
 - Supported by Intel and AMD (already OpenCL 2.0)
- •NVIDIA: Hardware support available, so just use PTX intrinsics
- •ARM: Use memory fences and hope for the best
 - Verify correctness using application test suites

Portable subgroups (OCL2.1)

- •Required for:
 - coop-conv,
 - nested parallelism (sg)
- •NVIDIA: Use PTX intrinsics
- Intel and AMD: Use vendor-specific extensions
 - Intel subgroup sizes can vary per kernel!
- •ARM: Assume subgroup size of 1

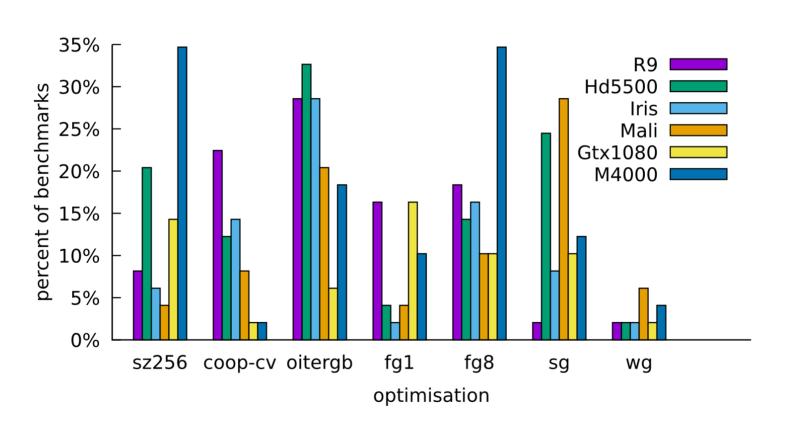
Portable global barriers (OCL-FP)

- •Required for: Iteration Outlining (oitergb)
- Assume occupancy-bound architecture
- •Kernel includes a prologue to detect how many workgroups are running (can undercount)
 - All workgroups with ID greater than this count exit
- Details in: Sorensen et al., "Portable interworkgroup barrier synchronisation for GPUs"

Experimental Setup

- .17 graph algorithms
- -3 graph inputs: Road, Random, R-MAT
- •6 GPUs, from 4 vendors (NVIDIA, Intel, ARM and AMD)

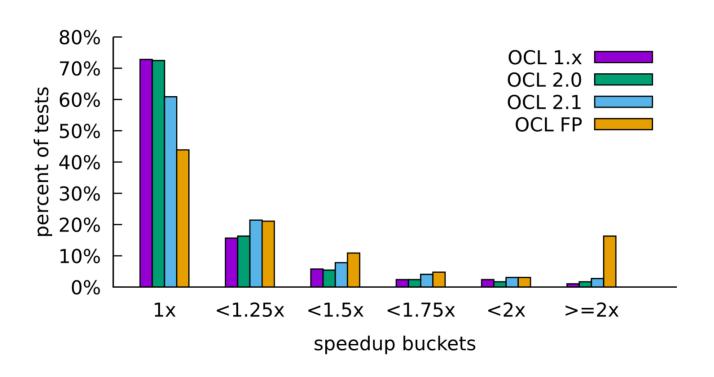
Portability of Optimisations



Portability of Optimisations (contd.)

- •All optimisations required to achieve speedup for some fraction of architecture + benchmark + input combinations
- Optimisations are not NVIDIA-specific

OpenCL Speedup Classes



Conclusion

- •Newer OpenCL features lead to better performance
- •Best performance obtained when OCL-FP features are used though these are not yet supported

Thank you!

•For more details, see Tyler's PhD Thesis

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